



COMPRESSED SENSING AND RELATED OPTIMIZATIONS

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OUTLINE

- CS overview
- Preliminaries
- Sparse recovery
- Optimization techniques
- Mixed-norm problems

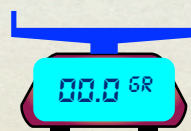


DIGITAL SCALE EXAMPLE

- 8 coins, one is fake



- Find the fake with a digital balance



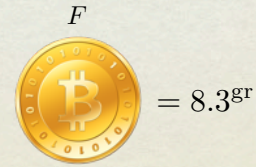


DIGITAL SCALE EXAMPLE

- 8 coins, one is fake



- None of the coins in experiment 1 is fake.
- All coins not used in experiment 2 are authentic.
- All coins not used in experiment 3 are authentic.



DIGITAL SCALE EXAMPLE

- Mathematical Formulation

$$\begin{bmatrix} 40 \\ 38.3 \\ 38.3 \end{bmatrix} = \underbrace{\begin{bmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 \end{bmatrix}}_{\Phi_{3 \times 8}} \begin{bmatrix} w_A \\ w_B \\ w_C \\ w_D \\ w_E \\ w_F \\ w_G \\ w_H \end{bmatrix}$$

$$\Rightarrow \underbrace{\begin{bmatrix} 0 \\ -1.7 \\ -1.7 \end{bmatrix}}_{\mathbf{y}_{3 \times 1}} = \Phi_{3 \times 8} \underbrace{\begin{bmatrix} w_A - 10 \\ w_B - 10 \\ w_C - 10 \\ w_D - 10 \\ w_E - 10 \\ w_F - 10 \\ w_G - 10 \\ w_H - 10 \end{bmatrix}}_{\mathbf{x}_{8 \times 1}}$$

$$\Rightarrow \mathbf{y}_{3 \times 1} = \Phi_{3 \times 8} \cdot \mathbf{x}_{8 \times 1}$$

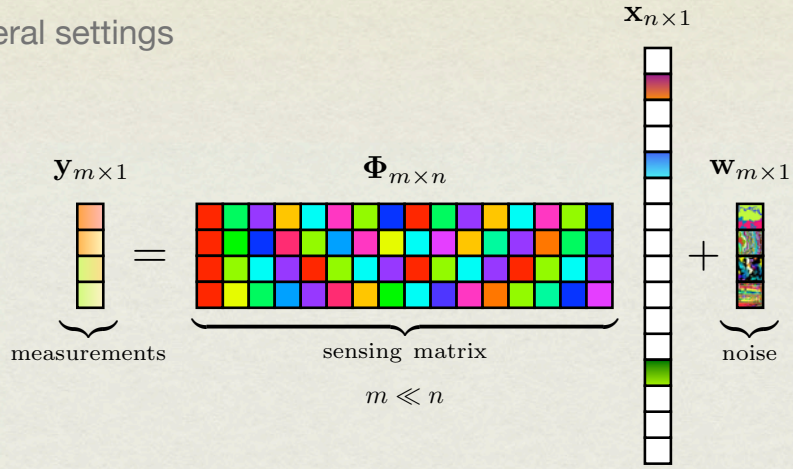
> Underdetermined LSoE!

> $\mathbf{x}_{8 \times 1} = 1$ -sparse!



CS PROBLEM

- More general settings



$$\begin{cases} y_{m \times 1} = \Phi_{m \times n} \cdot x_{n \times 1} + w_{m \times 1} \\ x_{n \times 1} = k\text{-sparse} \end{cases} \Rightarrow x_{n \times 1} = ?$$

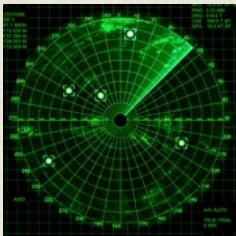
$m < n \Rightarrow$ infinitely many solutions



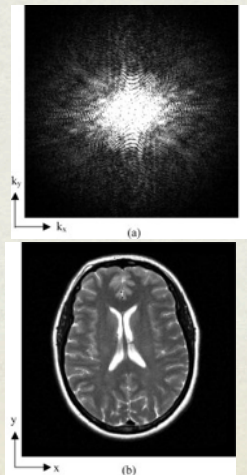
CS PROBLEM

- Similar problems

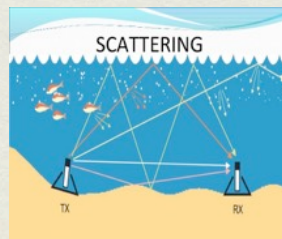
DOA estimation



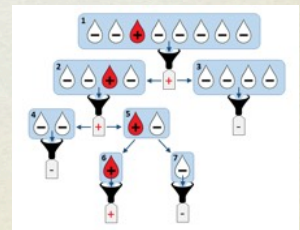
MRI scanning



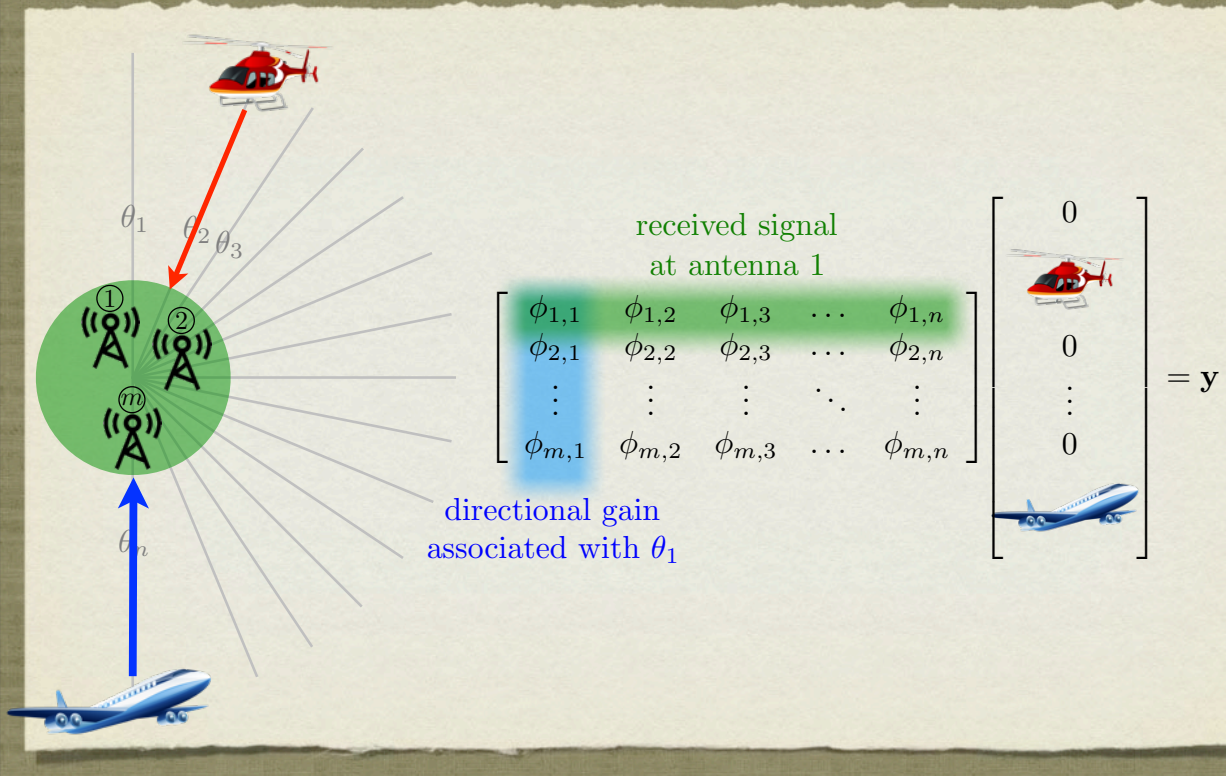
Channel estimation



Group testing



RADAR EXAMPLE





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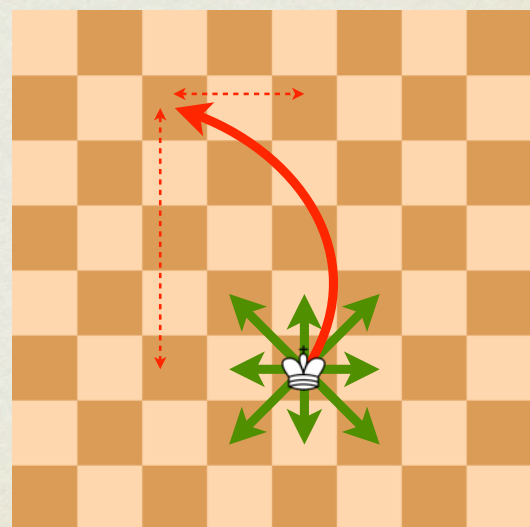
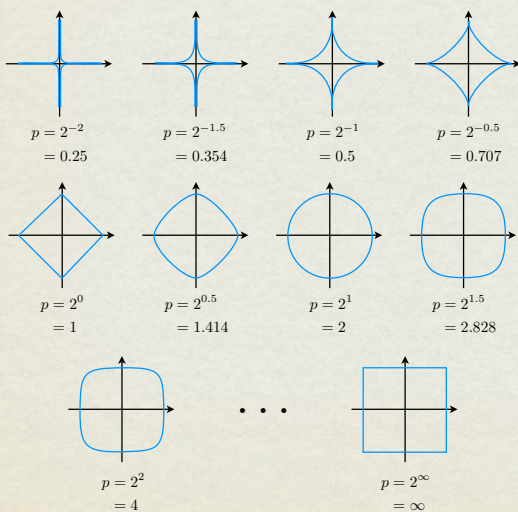


ℓ_p -NORMS

- Unit ball

$$\|\mathbf{x}_{n \times 1}\|_p = \sqrt[p]{|x_1|^p + \dots + |x_n|^p}$$

$$B_1(\mathbb{F}^n; \ell_p) \triangleq \{\mathbf{x} \in \mathbb{F}^n \mid \|\mathbf{x}\|_p = 1\}$$



SUITABLE SAMPLING



$$\text{SPARK}(\Phi_{m \times n}) \triangleq \min \left\{ t \mid \left\{ \begin{array}{l} \exists i_1, \dots, i_t \in \{1, \dots, n\} \\ \exists 0 \neq c_1, \dots, c_t \end{array} : \sum_{j=1}^t c_j \varphi_{i_j} = \mathbf{0}_{m \times 1} \right\} \right\}$$

$\Phi_{m \times n}$ satisfies $\text{NSP}(k, c_{\text{NSP}})$ iff

$$\forall \mathbf{h} \in \mathcal{N}(\Phi), \forall \Lambda \subset \{1, \dots, n\}, |\Lambda| = k : \|\mathbf{h}_\Lambda\|_2 < \frac{c_{\text{NSP}}}{\sqrt{k}} \|\mathbf{h}_{\Lambda^c}\|_1$$

$\Phi_{m \times n}$ satisfies $\text{RIP}(k, \delta_k)$ iff

$$0 \leq \delta_k < 1, \quad \forall \mathbf{x}_{n \times 1} \in \Sigma_k : (1 - \delta_k) \|\mathbf{x}\|_2^2 \leq \|\Phi_{m \times n} \mathbf{x}_{n \times 1}\|_2^2 \leq (1 + \delta_k) \|\mathbf{x}\|_2^2$$



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SPARSE RECOVERY

- How to recover \mathbf{x} ?

$$\mathbf{x}_{n \times 1} \in \Sigma_k, \quad \mathbf{y}_{m \times 1} = \Phi_{m \times n} \mathbf{x}_{n \times 1}$$

$$\mathbf{y}_{m \times 1} \stackrel{?}{\Rightarrow} \mathbf{x}_{n \times 1}$$

\mathbb{P}_0

$$\hat{\mathbf{x}}_{n \times 1} = \text{the sparsest } \mathbf{z}_{n \times 1}$$

$$\text{s.t. } \mathbf{y}_{m \times 1} = \Phi_{m \times n} \mathbf{z}_{n \times 1}$$

\equiv

$$\hat{\mathbf{x}}_{n \times 1} = \underset{\mathbf{z}_{n \times 1}}{\operatorname{argmin}} \|\mathbf{z}_{n \times 1}\|_0$$

$$\text{s.t. } \mathbf{y}_{m \times 1} = \Phi_{m \times n} \mathbf{z}_{n \times 1}$$

- $\text{SPARK}(\Phi) > 2k \Rightarrow \mathbf{x}_{n \times 1}$ is the sparsest solution to $\mathbb{P}_0 \Rightarrow \hat{\mathbf{x}}_{n \times 1} = \mathbf{x}_{n \times 1}$
- \mathbb{P}_0 is non-convex \Rightarrow NP-hard
- $\|\cdot\|_0$ is sensitive to noise

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SPARSE RECOVERY

- How to recover \mathbf{x} ?

$$\mathbf{x}_{n \times 1} \in \Sigma_k, \quad \mathbf{y}_{m \times 1} = \Phi_{m \times n} \mathbf{x}_{n \times 1}$$

$$\mathbf{y}_{m \times 1} \stackrel{?}{\Rightarrow} \mathbf{x}_{n \times 1}$$

- Convex relaxation

 \mathbb{P}_1

$$\begin{aligned} \hat{\mathbf{x}}_{n \times 1} = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} \|\mathbf{z}_{n \times 1}\|_1 \\ \text{s.t. } \mathbf{y}_{m \times 1} = \Phi_{m \times n} \mathbf{z}_{n \times 1} \end{aligned}$$

compressibility
instead
of sparsity

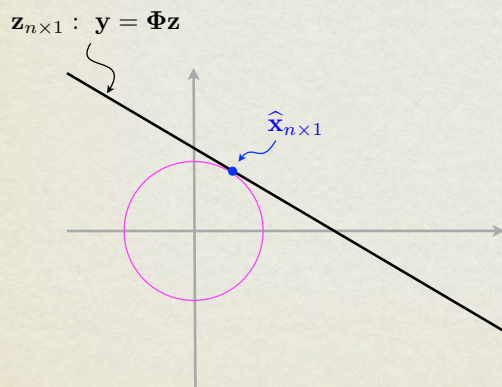
- \mathbb{P}_1 is convex \Rightarrow polynomial-time algorithm
- \mathbb{P}_1 does not necessarily result in a sparse solution
- $\Phi : \text{NSP}(k, 1) \Rightarrow \mathbf{x}_{n \times 1}$ is the unique minimizer of $\mathbb{P}_1 \Rightarrow \hat{\mathbf{x}}_{n \times 1} = \mathbf{x}_{n \times 1}$

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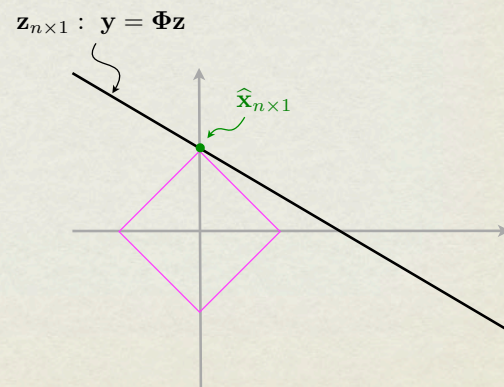


GEOMETRICAL PERSPECTIVE

$$\begin{aligned} \hat{\mathbf{x}}_{n \times 1} = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} \|\mathbf{z}_{n \times 1}\|_2 \\ \text{s.t. } \mathbf{y}_{m \times 1} = \Phi_{m \times n} \mathbf{z}_{n \times 1} \end{aligned}$$



$$\begin{aligned} \hat{\mathbf{x}}_{n \times 1} = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} \|\mathbf{z}_{n \times 1}\|_1 \\ \text{s.t. } \mathbf{y}_{m \times 1} = \Phi_{m \times n} \mathbf{z}_{n \times 1} \end{aligned}$$



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\mathbb{P}_1 VIA LP

\mathbb{P}_1

$$\begin{aligned} \widehat{\mathbf{x}}_{n \times 1} &= \underset{\mathbf{z}_{n \times 1}}{\operatorname{argmin}} \|\mathbf{z}_{n \times 1}\|_1 \\ \text{s.t. } \mathbf{y}_{m \times 1} &= \Phi_{m \times n} \mathbf{z}_{n \times 1} \end{aligned}$$

LP

$$\begin{aligned} \widehat{\mathbf{x}}_{n \times 1} &= \underset{\mathbf{z}_{n \times 1}}{\operatorname{argmax}} (\mathbf{c}_{n \times 1})^T \mathbf{z}_{n \times 1} \\ \text{s.t. } \begin{cases} \mathbf{z}_{n \times 1} & \geq \mathbf{0}_{n \times 1} \\ \mathbf{A}_{m \times n} \mathbf{z}_{n \times 1} & \leq \mathbf{b}_{m \times 1} \end{cases} \end{aligned}$$

\mathbb{P}_1 via LP

$$\begin{aligned} \widehat{\mathbf{x}}^{\pm}_{2n \times 1} &= \underset{\mathbf{z}_{2n \times 1}}{\operatorname{argmax}} [1, 1, \dots, 1]_{1 \times 2n} \mathbf{z}_{2n \times 1} \\ \text{s.t. } \begin{cases} \mathbf{z}_{2n \times 1} & \geq \mathbf{0}_{n \times 1} \\ \begin{bmatrix} \Phi & -\Phi \\ -\Phi & \Phi \end{bmatrix} \mathbf{z}_{2n \times 1} & \leq \begin{bmatrix} \mathbf{y} \\ -\mathbf{y} \end{bmatrix} \end{cases} \end{aligned} \quad \begin{aligned} \widehat{x}_i &= (\widehat{x}^{\pm})_i - (\widehat{x}^{\pm})_{n+i} \\ 1 &\leq i \leq n \end{aligned}$$



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LP SHORTCOMING

- Basis Pursuit (BP)

$$\mathbf{y}_{m \times n} = \Phi_{m \times n} \mathbf{x}_{n \times 1} \Rightarrow$$

\mathbb{P}_1

$$\begin{aligned} \hat{\mathbf{x}}_{n \times 1} = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} \|\mathbf{z}_{n \times 1}\|_1 \\ \text{s.t. } \mathbf{y}_{m \times 1} = \Phi_{m \times n} \mathbf{z}_{n \times 1} \end{aligned}$$

- Basis Pursuit DeNoising (BPDN)

$$\begin{aligned} \mathbf{y}_{m \times n} = \Phi_{m \times n} \mathbf{x}_{n \times 1} + \mathbf{w}_{m \times 1} \\ \|\mathbf{w}_{m \times 1}\|_2 \leq \epsilon \end{aligned} \Rightarrow$$

$\tilde{\mathbb{P}}_1$

$$\begin{aligned} \hat{\mathbf{x}}_{n \times 1} = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} \|\mathbf{z}_{n \times 1}\|_1 \\ \text{s.t. } \|\mathbf{y}_{m \times 1} - \Phi_{m \times n} \mathbf{z}_{n \times 1}\|_2 \leq \epsilon \end{aligned}$$

- Dantzig Selector

$$\begin{aligned} \mathbf{y}_{m \times n} = \Phi_{m \times n} \mathbf{x}_{n \times 1} + \mathbf{w}_{m \times 1} \\ \|\Phi_{m \times n}^T \mathbf{w}_{m \times 1}\|_\infty \leq \omega \end{aligned} \Rightarrow$$

DS

$$\begin{aligned} \hat{\mathbf{x}}_{n \times 1} = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} \|\mathbf{z}_{n \times 1}\|_1 \\ \text{s.t. } \|\Phi^T (\mathbf{y} - \Phi \mathbf{z})\|_\infty \leq \omega \end{aligned}$$

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LAGRANGE FORM

- BP / BPDN

 \mathbb{P}_1

$$\begin{aligned} \hat{\mathbf{x}}_{n \times 1} = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} \|\mathbf{z}_{n \times 1}\|_1 \\ \text{s.t. } \mathbf{y}_{m \times 1} = \Phi_{m \times n} \mathbf{z}_{n \times 1} \end{aligned}$$

 $\tilde{\mathbb{P}}_1$

$$\begin{aligned} \hat{\mathbf{x}}_{n \times 1} = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} \|\mathbf{z}_{n \times 1}\|_1 \\ \text{s.t. } \|\mathbf{y}_{m \times 1} - \Phi_{m \times n} \mathbf{z}_{n \times 1}\|_2 \leq \epsilon \end{aligned}$$

- Lagrangian form

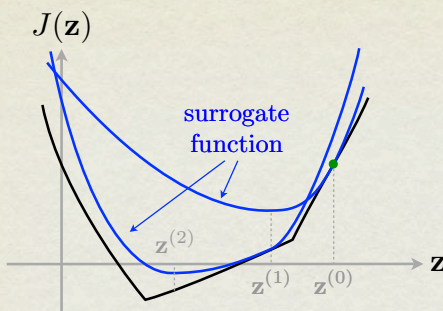
$$\hat{\mathbf{x}}_{n \times 1} = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} J(\mathbf{z}_{n \times 1}) = \operatorname{argmin}_{\mathbf{z}_{n \times 1}} \|\mathbf{y}_{m \times 1} - \Phi_{m \times n} \mathbf{z}_{n \times 1}\|_2^2 + \lambda \|\mathbf{z}_{n \times 1}\|_1$$

- λ should be properly set.
- $J(\mathbf{z}_{n \times 1})$ is convex.
- $J(\mathbf{z}_{n \times 1})$ is non-smooth.

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SURROGATE FUNCTIONS



- Surrogate function $\tilde{J}(\mathbf{z}_{n \times 1})$:
 - $\forall \mathbf{z}_{n \times 1} : \tilde{J}(\mathbf{z}) \geq J(\mathbf{z})$.
 - For a given \mathbf{z}^* : $\tilde{J}(\mathbf{z}^*) = J(\mathbf{z}^*)$.
 - Minimizing \tilde{J} is easy.

$$\begin{cases} J(\mathbf{z}^{(0)}) = \tilde{J}(\mathbf{z}^{(0)}) \\ \tilde{J}(\mathbf{z}^{(1)}) \leq \tilde{J}(\mathbf{z}^{(0)}) \\ J(\mathbf{z}^{(1)}) \leq \tilde{J}(\mathbf{z}^{(1)}) \end{cases} \Rightarrow J(\mathbf{z}^{(1)}) \leq J(\mathbf{z}^{(0)}) \Rightarrow \dots$$

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IRLS

$$J(\mathbf{z}_{n \times 1}) = \|\mathbf{y}_{m \times 1} - \Phi_{m \times n} \mathbf{z}_{n \times 1}\|_2^2 + \lambda \|\mathbf{z}_{n \times 1}\|_1$$

useful inequality: $\forall a, b, |a| \leq \frac{1}{2} \frac{a^2}{|b|} + \frac{1}{2}|b|$

↑
equality for
 $a = b$

$$\tilde{J}_{i+1}(\mathbf{z}_{n \times 1}) = \|\mathbf{y}_{m \times 1} - \Phi_{m \times n} \mathbf{z}_{n \times 1}\|_2^2 + \frac{\lambda}{2} \sum_{j=1}^n \left(\frac{z_j^2}{|z_j^{(i)}|} + |z_j^{(i)}| \right)$$

previous iteration

IRLS: $\left\{ \begin{array}{l} \text{initialize } \mathbf{z}^{(0)} \\ \text{set } i = 0 \\ \text{while (stopping condition)} \\ \quad \mathbf{z}^{(i+1)} = \left(\Phi^H \Phi + \frac{\lambda}{2} \text{diag} \left(\frac{1}{|z_1^{(i)}| + \epsilon}, \dots, \frac{1}{|z_n^{(i)}| + \epsilon} \right) \right)^{-1} \Phi^H \mathbf{y} \\ \quad i \leftarrow i + 1 \\ \quad \text{update stop. cond.} \\ \text{end} \end{array} \right.$



ISTA

$$J(\mathbf{z}_{n \times 1}) = \|\mathbf{y}_{m \times 1} - \Phi_{m \times n} \mathbf{z}_{n \times 1}\|_2^2 + \lambda \|\mathbf{z}_{n \times 1}\|_1$$

useful inequality: $\forall \mathbf{v}_{n \times 1}, \|\Phi_{m \times n} \mathbf{v}_{n \times 1}\|_2^2 \leq \underbrace{\lambda_{\max}(\Phi^H \Phi)}_L \|\mathbf{v}\|_2^2$

$$\tilde{J}_{i+1}(\mathbf{z}_{n \times 1}) = \|\mathbf{y} - \Phi \mathbf{z}\|_2^2 + L \|\mathbf{z} - \mathbf{z}^{(i)}\|_2^2 - \|\Phi(\mathbf{z} - \mathbf{z}^{(i)})\|_2^2 + \lambda \|\mathbf{z}\|_1$$

previous iteration

$$= L \|\mathbf{z} - \left(\mathbf{z}^{(i)} - \frac{1}{L} \Phi^H (\Phi \mathbf{z}^{(i)} - \mathbf{y}) \right)\|_2^2 + \lambda \|\mathbf{z}\|_1 + \left(\begin{array}{l} \text{terms not} \\ \text{having } \mathbf{z} \end{array} \right)$$

$$\underset{\mathbf{z}_{n \times 1}}{\text{argmin}} \tilde{J}_{i+1}(\mathbf{z}) \Rightarrow \underset{z_j}{\text{argmin}} |z_j - \left(\mathbf{z}^{(i)} - \frac{1}{L} \Phi^H (\Phi \mathbf{z}^{(i)} - \mathbf{y}) \right)_j|^2 + \frac{\lambda}{L} |z_j|$$

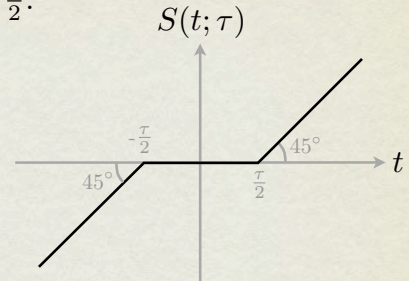
$$\mathbb{F}^n \text{ minimisation} \equiv n \times (\mathbb{F}^1 \text{ minimisation})$$



ISTA

$$S(t; \tau) = \underset{z}{\operatorname{argmin}} |z - t|^2 + \tau |z| = \begin{cases} (|t| - \frac{\tau}{2}) e^{j\angle t}, & |t| > \frac{\tau}{2}, \\ 0, & |t| \leq \frac{\tau}{2}. \end{cases}$$

shrinkage
function



ISTA: $\left\{ \begin{array}{l} \text{initialize } \mathbf{z}^{(0)} \\ \text{set } L = \lambda_{\max}(\Phi^H \Phi) \\ \text{set } i = 0 \\ \text{while (stopping condition)} \\ \quad \mathbf{t} \leftarrow \mathbf{z}^{(i)} - \frac{1}{L} \Phi^H (\Phi \mathbf{z}^{(i)} - \mathbf{y}) \\ \quad \mathbf{z}^{(i+1)} = [S(t_1; \frac{\lambda}{L}), \dots, S(t_n; \frac{\lambda}{L})]^T \\ \quad i \leftarrow i + 1 \\ \quad \text{update stop. cond.} \\ \text{end} \end{array} \right.$

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ADMM

$$\mathbf{x}_{n \times 1}^* = \underset{\mathbf{z}_{n \times 1}}{\operatorname{argmin}} f(\mathbf{z}) + g(\mathbf{z})$$

• ADMM

$$\begin{aligned} \mathbf{x}_{n \times 1}^* &= \underset{\mathbf{z}_{n \times 1}, \zeta_{n \times 1}}{\operatorname{argmin}} f(\mathbf{z}) + g(\zeta) \\ \text{s.t. } &\mathbf{z} - \zeta = 0 \end{aligned}$$

$$J_k(\mathbf{z}, \zeta) \triangleq f(\mathbf{z}) + g(\zeta) + \mu \|\mathbf{z} - \zeta\|_2^2 - (\tau_k)^T (\mathbf{z} - \zeta)$$

$$\left\{ \begin{array}{l} \tau_{k+1} = \tau_k - 2\mu(\mathbf{z}^{(k)} - \zeta^{(k)}) \\ \mathbf{z}^{(k+1)} = \underset{\mathbf{z}_{n \times 1}}{\operatorname{argmin}} J_{k+1}(\mathbf{z}, \zeta^{(k)}) \\ \zeta^{(k+1)} = \underset{\zeta_{n \times 1}}{\operatorname{argmin}} J_{k+1}(\mathbf{z}^{(k+1)}, \zeta) \end{array} \right.$$

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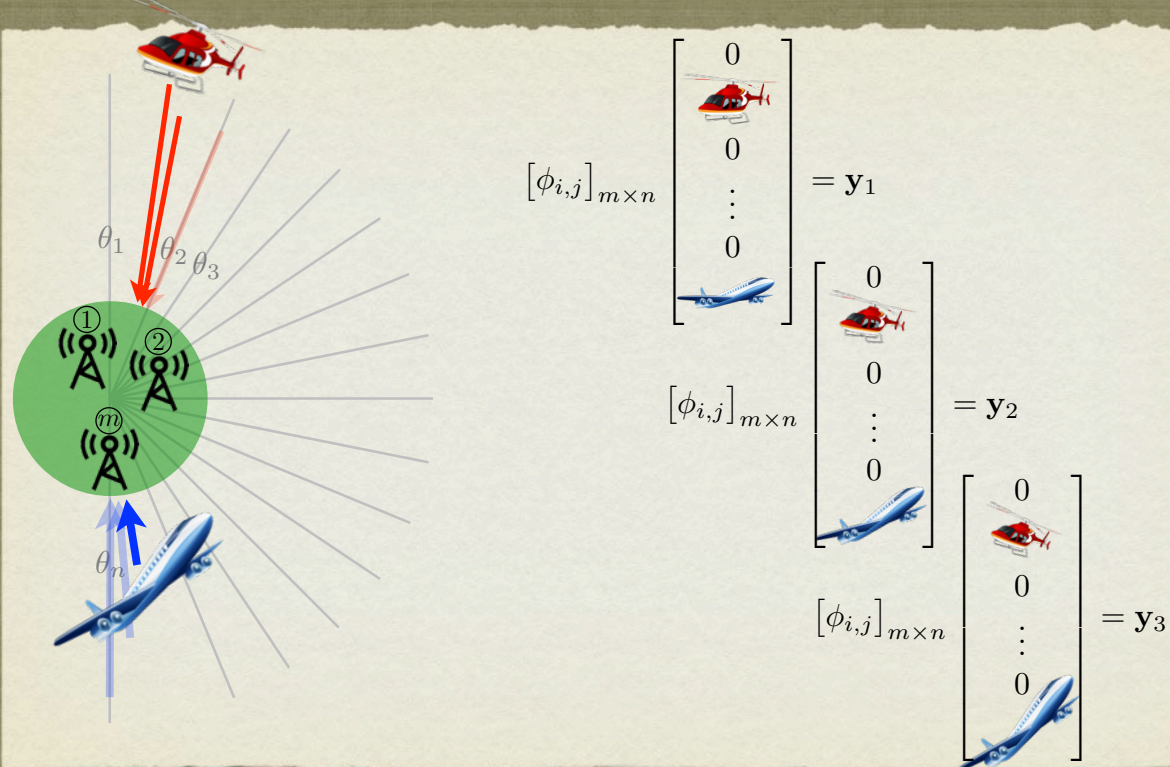


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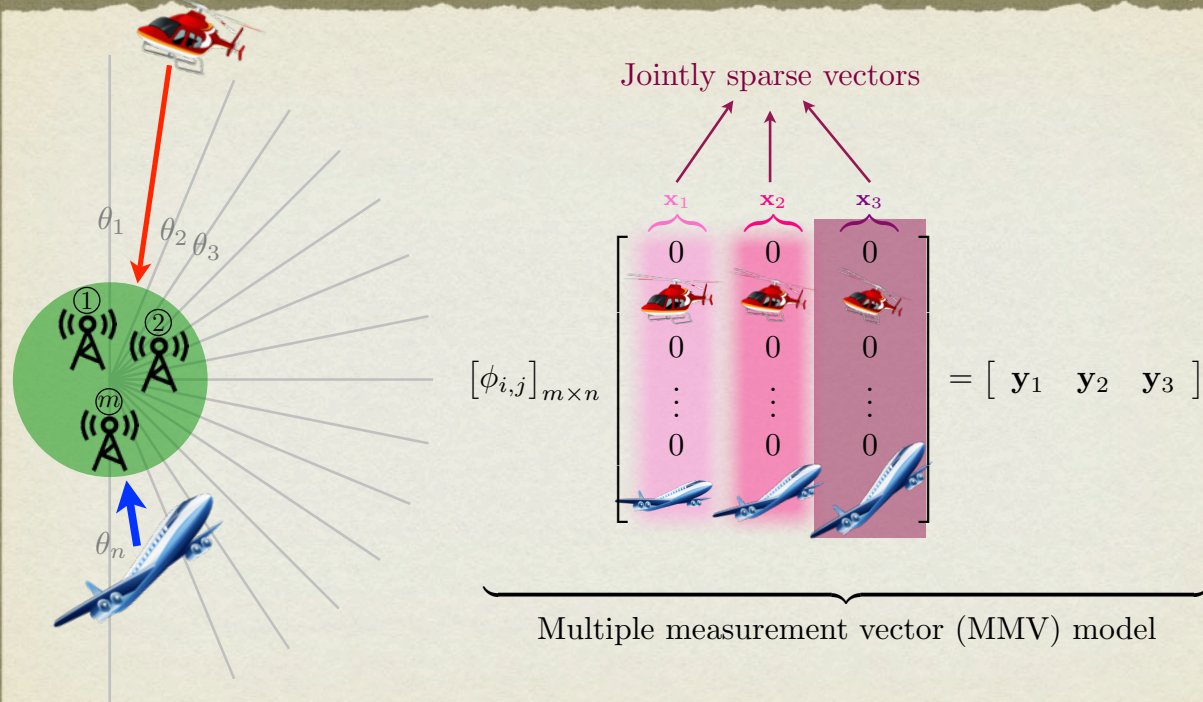
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MULTI-SNAPSHOT RADAR



MULTI-SNAPSHOT RADAR



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MMV



$$\left\{ \begin{array}{l} \underbrace{\mathbf{Y}_{m \times d}}_{\text{measurements}} = \underbrace{\Phi_{m \times n}}_{\text{sensing matrix}} \underbrace{\mathbf{X}_{n \times d}}_{\text{jointly sparse}} \\ \mathbf{X}_{n \times d} \text{ has few non-zero rows} \end{array} \right. \quad \text{or} \quad \left\{ \begin{array}{l} \underbrace{\mathbf{Y}_{m \times d}}_{\text{measurements}} = \underbrace{\Phi_{m \times n}}_{\text{sensing matrix}} \underbrace{\mathbf{X}_{n \times d}}_{\text{jointly sparse}} + \underbrace{\mathbf{W}_{m \times d}}_{\text{noise}} \\ \mathbf{X}_{n \times d} \text{ has few non-zero rows} \end{array} \right.$$

$$\hat{\mathbf{X}}_{n \times d} = \underset{\mathbf{Z}_{n \times d}}{\operatorname{argmin}} \|\mathbf{Z}_{n \times d}\|_{2,1}$$

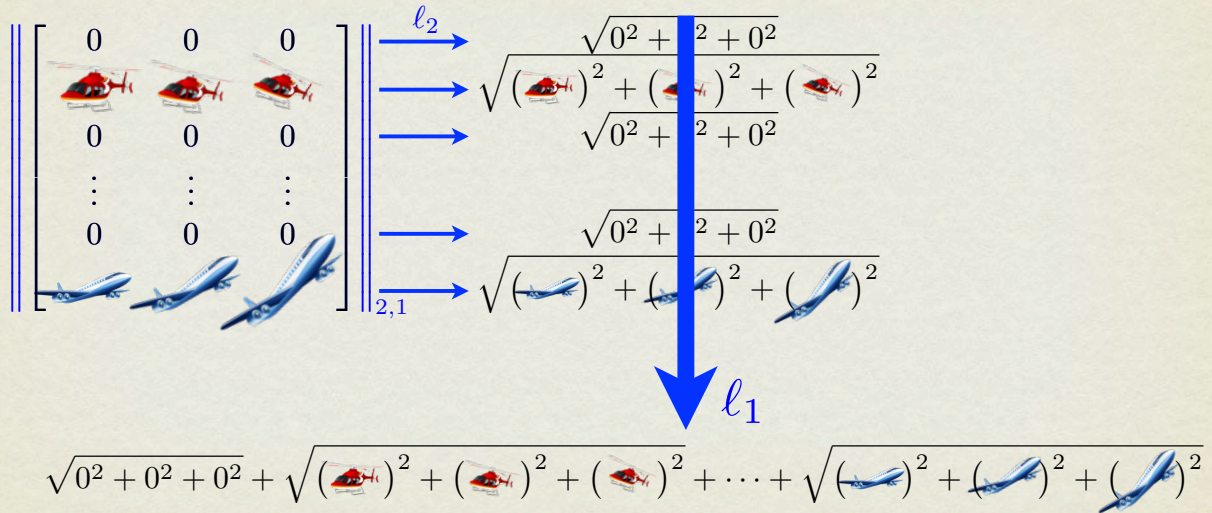
s.t. $\mathbf{Y}_{m \times d} = \Phi_{m \times n} \mathbf{Z}_{n \times d}$

$$\hat{\mathbf{X}}_{n \times d} = \underset{\mathbf{Z}_{n \times d}}{\operatorname{argmin}} \|\mathbf{Z}_{n \times d}\|_{2,1}$$

s.t. $\|\mathbf{Y}_{m \times d} - \Phi_{m \times n} \mathbf{Z}_{n \times d}\|_2 \leq \epsilon$

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$\ell_{2,1}$



$\ell_{2,1}$

$$\hat{\mathbf{X}}_{n \times d} = \underset{\mathbf{Z}_{n \times d}}{\operatorname{argmin}} \|\mathbf{Z}_{n \times d}\|_{2,1}$$

s.t. $\mathbf{Y}_{m \times d} = \Phi_{m \times n} \mathbf{Z}_{n \times d}$

Second order
cone programming
(SOCP)

$$\underset{t_1, \dots, t_n, \mathbf{Z}_{n \times d}}{\operatorname{argmin}} \sum_{i=1}^n t_i$$

s.t. $\begin{cases} 0 \leq t_i & , \forall 1 \leq i \leq n \\ \|[Z_{i,1}, Z_{i,2}, \dots, Z_{i,d}]\|_2 \leq t_i & , \forall 1 \leq i \leq n \\ \mathbf{Y}_{m \times d} = \Phi_{m \times n} \mathbf{Z}_{n \times d} \end{cases}$

Lagrange form

+
ADMM

$$\underset{\mathbf{Z}_{n \times d}, \mathbf{B}_{m \times d}, \mathbf{C}_{m \times d}}{\operatorname{argmin}} \|\mathbf{Y} - \underbrace{\mathbf{B}}_{\Phi \mathbf{Z}}\|_2^2 + \lambda \|\mathbf{Z}\|_{2,1} + \mu \|\mathbf{B} - \Phi \mathbf{Z}\|_2^2 - \operatorname{tr}(\mathbf{C}^T (\mathbf{B} - \Phi \mathbf{Z}))$$

Thank you!



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